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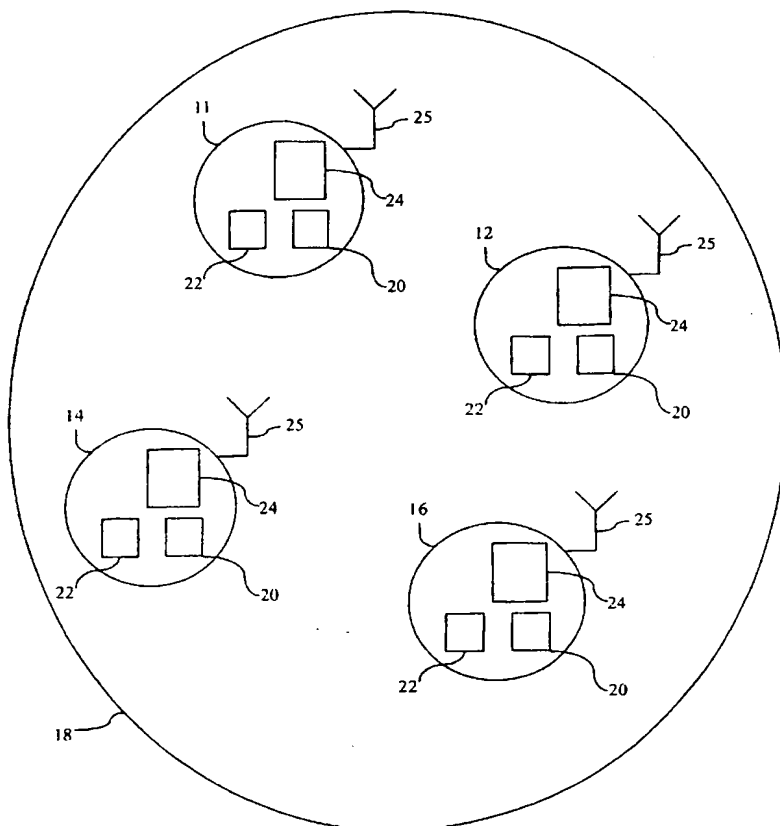
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(54) Title: LOCAL AREA NETWORK



(57) Abstract: A method and system for distributed security for a plurality of devices in a communication network, each of the devices being responsible for generating, distributing and controlling its own keys for access to the communication network and using the keys to establish a trusted network, each device's membership to the communication network being checked periodically by other devices by using a challenge response protocol to establish which devices are allowed access to the communication network and the trusted network.

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1 LOCAL AREA NETWORK

2 BACKGROUND OF THE INVENTION

3 [0001] This application claims priority in United States Provisional Application
4
5 Serial No. 60/362,865, entitled "Local Area Network", filed on March 8, 2002 and
6
7 United States Provisional Application Serial No. 60/363,309, entitled "Local Area
8
9 Network", filed on March 11, 2002.

10 FIELD OF THE INVENTION

11 [0002] This invention relates to communication networks, more particularly it relates
12 to security within these networks.
13

14 DESCRIPTION OF THE PRIOR ART

15
16 [0003] One of the most significant recent developments in wireless technologies is
17 the emergence of wireless personal area networking. Wireless personal area networks
18 WPANs™ use radio frequencies to transmit both voice and data, and are specified by
19 standards such as IEEE standard 802.15 or 802.3 from the Institute of Electrical and
20 Electronics Engineers Standards Association (IEEE-SA), among other specifications. The
21 802.15 specification is ideal for linking notebook computers, mobile phones, personal
22 digital assistants (PDAs), digital cameras, and other handheld devices to do business at
23 home, on the road, or in the office.

24 [0004] These wireless networks are formed by a number of devices joining and
25 leaving the network in an ad hoc manner, hence such networks are known as ad hoc
26 networks or piconets. Thus, the set of devices connected to the ad hoc network any given
27 time may fluctuate, and so the topology of the network is dynamic. It is desirable to
28 control access to the network and to provide a mechanism for establishing and
29 maintaining security. Traditionally, security is established using a central device or a
30 piconet controller (PNC) which controls access and distributes keys within the network.
31 A drawback of this scheme is that each member of the network is required to trust the
32 PNC.

1 [0005] Admission to the piconet is based on the outcome of the following protocols
2 between the prospective joining device and the PNC of the piconet. The joining device
3 and the PNC engage in a mutual entity authentication protocol based on public key or
4 symmetric key techniques. The true device identity of both the joining device and the
5 PNC is determined using this protocol. A link key can also be derived based on the
6 authentic keys of both parties. Another protocol involves using authorization techniques
7 between both devices, based on access control lists (ACLs). The Access Control Lists
8 may be dynamically updated, similar to PDA functionality, where a determination is
9 made whether an entity is added or removed from the ACL at entry. This determination
10 may be made by an operator, such as a human operator. For devices that lack a user
11 interface, this update mechanism may be invoked by an open enrollment period followed
12 by a lock-up step, for example, which may be confirmed by a button push or be a simple
13 re-set of the whole list. This may be performed by actuating a re-set or re-initialize button
14 on the device.

15 [0006] Thus devices in the piconet fully depend on information provided by the PNC
16 regarding which devices have been admitted to the piconet, since admission is based on
17 communication between the PNC and a joining device only. If however an improper list
18 of devices, DeviceList, in the piconet has been distributed by the PNC, either by error or
19 maliciously, the security of the network is jeopardised. Each device has a short hand
20 address, such as a local 8-bit ID, and a long hand address, such as a global 48-bit device
21 ID. For example, in a piconet in which since all devices share a common broadcast key,
22 the list of admitted devices to the piconet is $L := (\text{local 8-bit device ID, global 48-bit}$
23 $\text{device ID})$, then the failure to obtain the complete and authentic list of admitted devices
24 has the following consequences:

25 [0007] *'Fly on the wall' scenario:*

26 [0008] If a device obtains an incomplete list: $L' \subset (L' \neq L)$ of admitted devices, all
27 devices in the complementary set $L \setminus L'$ are 'invisible' to the device. Hence, the device
28 might mistakenly think it is sharing secured information only with devices from the list
29 L' , whereas actually it is unknowingly sharing with other devices of the set L as well.
30 This obviously violates sound security practice.

31 [0009] *'Switchboard' scenario ':*

1 **[0010]** If the binding between the local device ID and the global device ID is
2 incorrectly received, for example if 2 entries are interchanged, a device might direct
3 information to the improper device and so compromise the intended security. This
4 property also holds in other settings where a key-generating party does not share
5 complete and authentic information on the composition of the key-sharing group itself
6 with the other members of this group. Therefore, these scenarios present a security model
7 in which there is complete trust or a security model in which a device trusts no other
8 device, however a hybrid model of these two models is possible.

9 **[0011]** Accordingly it is an object of the present invention to mitigate or obviate at
10 least one of above-mentioned disadvantages.

11 12 SUMMARY OF THE INVENTION

13
14 **[0012]** In one of its aspects the invention provides a method of establishing and
15 maintaining distributed security between a plurality of devices in an ad hoc network, the
16 method having the steps of; associating each device with a unique device address;
17 assigning to one of the devices a control function to control access to the network
18 by other devices;
19 each of the devices generating a public key for distribution to other devices;
20 each of the devices authenticating itself periodically with the other devices in order to
21 determine status of the other devices;
22 arranging the devices into a plurality of trust groups, each group having a group
23 key for distribution within the trust group;
24 associating a trust level to each of the devices;
25 each of the devices using the public key and the group key to perform key
26 agreement in order to establish a secure communication channel with the other devices in
27 the group;
28 whereby each of the devices is responsible for its own security by generating,
29 distributing its own keys to the other devices.

1 **[0013]** In another aspect, the invention provides a method of establishing and
2 maintaining distributed security between one correspondent and another correspondent,
3 the correspondents being members of different ad hoc networks and forming a group of
4 communicating correspondents, the method having the steps of;
5 associating the one correspondent and the other correspondent with unique device
6 addresses;
7 controlling access to the different ad hoc networks;
8 each ad hoc network having a gateway and transferring traffic between the
9 correspondents via the gateways;
10 the one correspondent generating a public key for distribution to the other
11 correspondent;
12 the one correspondent authenticating itself periodically with the other
13 correspondent in order to determine status of the other correspondent;
14 determining a group key for distribution to the correspondents in accordance to
15 the step of controlling access;
16 associating a trust level to each correspondent; each of the correspondents using
17 the public key and the group key for performing key agreement in order to establish
18 secure communication within the group;
19 whereby the one correspondent is responsible for its own security by generating,
20 distributing its own keys to the other correspondent.

21 **[0014]** In yet another aspect, the invention provides a distributed security system for a
22 plurality of devices in a network, each of the devices being responsible for generating,
23 distributing and controlling its own keys for access to the network and using the keys to
24 establish a trusted network, each device's membership to the network being checked
25 periodically by other devices by using a challenge response protocol to establish which
26 devices are allowed access to the network and the trusted network.

28 BRIEF DESCRIPTION OF THE DRAWINGS

29 **[0015]** These and other features of the preferred embodiments of the invention will
30 become more apparent in the following detailed description in which reference is made to
31 the appended drawings wherein

1 [0016] Figure 1 is a communication network;

2 [0017] Figure 2 is a group structure for a security model having different trust levels;

3 [0018] Figure 3 is a group structure for a security model having different trust levels;

4 [0019] Figure 4 is a group structure for a security model having different trust levels;

5 [0020] Figure 5 is a group structure for a security model having different trust levels;

6 [0021] Figure 6 shows communication between piconets;

7 [0022] Figure 7 shows a flowchart outlining steps for establishing secure

8 communication between devices in different piconets; and

9 [0023] Figure 8 shows secure communication between piconets;

11 DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

12 [0024] Reference is first made to Figure 1, which shows an overview of a distributed
13 security system 10 having a plurality of communication devices 11, 12, 14, 16 in a
14 communication network 18, in a preferred embodiment. The communication network 18
15 may be a wireless personal area network (WPANTM) such as a piconet, in which the
16 devices 11, 12, 14, 16 connect to each other in an ad hoc fashion. The devices 11, 12, 14,
17 16 may be portable and mobile computing devices such as PCs, Personal Digital
18 Assistants (PDAs), peripherals, cell phones, pagers, consumer electronics, and other
19 handheld devices. It will be understood that such devices 11, 12, 14, 16 include
20 addressing information to facilitate communication within the network 18. The
21 addressing information includes a local device ID, having 8 bits for example, and a
22 device ID, such as, an IEEE MAC Address including 48 bits. Therefore, upon a device
23 11, 12, 14, 16 joining the network it is assigned an unused local ID. Generally, one device
24 11 will act as a master or a piconet network controller (PNC), and the other devices 12,
25 14, 16 act as slaves for the duration of the piconet 18 connection. The PNC 11 sets a
26 clock, a hopping pattern determined by device ID, and assigns time for connections
27 between all devices 11, 12, 14 16. Thus, each piconet 18 includes a unique hopping
28 pattern/ID, and the PNC 11 gives slaves 12, 14 16 the clock and a local device ID, which
29 is optionally used in conjunction with the IEEE MAC Address, to form the piconet 18.

30 [0025] The PNC 11 activates an access controller 20 using ID's of the devices and
31 optionally an access control list such that devices 12, 14, 16 that have been positively

1 authenticated and have been authorized are admitted to the piconet 18. The PNC 11 also
2 includes a traffic controller 22 to regulate data flow within the network 18. This may be
3 done by allocating time slots to each device 11,12,14,16 for message distribution. Each of
4 the devices 11, 12, 14, 16 includes a security manager function 24. The security manager
5 function 24 generates keys for communicating with other devices 11,12,14,16 within the
6 network 18, and distributes these keys to selected device members 11,12,14,16 of the
7 network 18. Each device 11, 12, 14 or 16 includes a transceiver 25 for establishing a
8 communication channel with other devices 11,12,14,16. When distributing a key, the
9 security manager function 24 also indicates to the other devices 11,12,14,16 in the
10 network 18 the other devices 11,12,14,16 to which the key is being distributed. Thus,
11 there is no reliance on other devices 11, 12, 14, 16 for trust functionality, as each device
12 11, 12, 14 or 16 need only trust itself, to form a distributed security regime.

13 [0026] Thus, the security manager function 24 can establish a trust set, or TrustList,
14 which indicates which of the devices 11,12,14,16 in the network the security manager 24
15 of that particular device 11,12,14 or 16 is prepared to trust. The security manager
16 function 24 may also attribute different levels of trust to each of the established trust sets.
17 In this way the equivalent of a centralised network 18 can be established where a device
18 11,12,14 or 16 trusts every other device 11,12,14 or 16; or an entirely decentralised
19 network 18 is provided where a device 11,12,14 or 16 trusts no other device 11,12,14 or
20 16 but itself.

21 [0027] Similarly the security manager 24 receiving a key from another device 11, 12,
22 14, 16 can determine its source and allocate to that key a level of trust that determines the
23 functions for which the key will be used. Thus the security manager 24 may determine
24 that the key is from a trusted party 11, 12, 14 or 16 and the key may be used to both
25 decrypt messages received from that trusted party 11, 12, 14 or 16 and encrypt messages
26 sent to that trusted party 11, 12, 14 or 16. Alternatively, the security manager function 24
27 may determine that the key originates at a party 11, 12, 14 or 16 not trusted by itself and
28 only permit the key to be used for decryption. However, the device 11, 12, 14 or 16 may
29 choose to ignore data, rather than going through the effort of having to decrypt the data
30 first. This option may be useful for dealing with unsolicited communication or
31 'junkmail'.

1 [0028] The security manager 24 also includes methods of determining which of the
2 devices 11, 12, 14 or 16 are presently active in the network 18. These methods include
3 the functions of each device 11, 12, 14 or 16 re-authenticating itself with each of its key
4 sharing parties 11, 12, 14 or 16 at predetermined time. One such method includes the
5 steps of periodically performing a 'heartbeat operation' in the form of a challenge
6 response protocol to determine which devices are presently included in the network 18,
7 and adjusting the groups and trust levels accordingly. Thus, each device 11, 12, 14 or 16
8 may dynamically update its own TrustList to reflect changes in the trust relationships. For
9 devices 11,12,14 or 16 that lack a user interface, this update mechanism may be invoked
10 by an open enrollment period followed by a lock-up step, possibly confirmed by a button
11 push, or it may be a simple re-set of the whole list, for example by pushing a re-set or re-
12 initialize button on the device 11,12,14 or 16.. Moreover, some of the changes might be
13 invoked by a third entity that performs remote or delegated trust management for that
14 device.

15 [0029] Referring now to Figure 2, in order to describe the distributed security model,
16 as an example, assume the PNC 11 permits access to devices A, B,C,D, E, F, G, H, then
17 the DeviceSet := {A,B,C,D,E,F,G,H}. However if the device A only trusts devices A, B,
18 C then TrustSet(A) := {A, B, C} that is Group 1. Also, device A may participate in other
19 groups having a different trust set, such as Group 2, having only device D. Thus the
20 security manger function 24 of device A senses Group 1 and Group 2 with different
21 constituent members and different levels of trust. For example, in Group 1, if device C is
22 the key source, and since device C is part of the TrustSet(A), this key by device C is
23 distributed which is used for both encryption/decryption permitted as C, and device A
24 only accepts keys transferred to itself by devices $DEV \in \text{TrustSet}(A)$, for encryption and
25 decryption purposes. In Group 2, as device D is not part of TrustSet(A), then A accepts a
26 key from device D, and any other devices E, F,G and H, which are not part of
27 TrustSet(A), for decryption purposes only. Accordingly if device A desires to
28 communicate to Group2 members, the device A generates a new group key to form a new
29 group, Group 3, and device A distributes this new group key to the members of Group2',
30 that is device D. Therefore, the groups then under the control of the security manager of

1 device A will then be Group 1, Group 2, as mentioned above, and Group 3, as shown
2 Figure 3.

3 **[0030]** The flexibility of the security managers 24 of devices A, B, C, D, E, F, G, H
4 permits different network structures to be mimicked. For example, using the notation
5 above, if $\text{DeviceSet} := \{A, B, C, D, E, F, G, H\}$, and $\text{TrustSet}(A) := \text{Universe}$, then device A
6 can be considered an altruistic device which provides a structure equivalent to a
7 centralized model. Conversely, if $\text{TrustSet}(D) := \{D\}$, then device D is an egocentric
8 device, and is a structure equivalent to a completely decentralized model. Then, looking at
9 Figure 4, device A participates in Groups 1, 2 and 3, all groups having with differing trust
10 relationships. For example, in Group 1 having devices A, B and C, if the key source is
11 device C, then this group key is used for encryption and decryption, as device A trusts all
12 devices B, C, D, E, F, G and H, which of course includes the key source C. However, in
13 Group 2 having devices A, D, and G, with the key source being device G, once again
14 device A uses this group key is used for encryption and decryption, while device D uses
15 it for decryption only as it does not trust any other device A, B, C, E, F, G or H. In Group 3
16 having devices D and E, with the key source being device E, device D uses the group key
17 for decryption only as it does not trust device E. As device A is not included in Group 3,
18 it does not receive the key.

19 **[0031]** In Figure 5, where one of the device F is hidden from the other members in
20 the network 18, then Group 2 does not include the full list of member devices, A, D, G and
21 H. Therefore, device D can not communicate with device F as the heartbeat operation
22 will indicate that device D is not alive. Since the 8-bit address or the 48-bit address of
23 device is unavailable, there is no communication between D and device F. Therefore,
24 device D uses the group keys for decryption only.

25 **[0032]** Thus, these different group structures as shown in Figures 2, 3, 4 and 5 may
26 be established within the same network 18 by using a decentralised or distributed security
27 management scheme having the ability to set different levels of trust per device. This may
28 be used in a number of ways, such as admission of devices A, B, C, D, E, F, G and H,
29 such as PDAs to a piconet 18 based on different subscription models. For example, one
30 subscription model may include charging a fee for airtime/bandwidth fee, while another
31 model may be based on charging for content. In this example, the models may be

1 implemented in a building, such as an airport or fitness club, the network 18 includes a
2 fixed PNC 11 on a ceiling and the PNC 11 multicasting to subscribing devices only, or
3 the models may be implemented between individual devices. Thus, by separating the role
4 of the security manager 24 from that of the PNC 11, charging models that differentiate
5 between airtime/bandwidth cost and content/subscription cost are possible, as these
6 charging models might be operated by different entities A,B,C,D,E,F,G or H, or another
7 intermediate entity.

8 [0033] It will be seen therefore that a versatile network 18 is provided, and moreover
9 the removal of a device A,B,C,D,E,F,G or H from the network 18 does not require re-
10 establishment of all keys in the network 18 as the individual devices A,B,C,D,E,F,G or H
11 control the distribution of the keys. Figure 6 shows communication between a device A in
12 piconet 1 with another device B in piconet 2, where Z_1 and Z_2 are members of piconet 1
13 and piconet 2, respectively. Z_1 and Z_2 include transceivers 25 for establishing a
14 communication channel or relay channel 26 between piconet 1 and piconet 2. Thus, Z_1
15 listens in on all traffic and sends all traffic destined for device B to Z_2 via the relay
16 channel 26. Upon receipt of the traffic relayed by Z_1 , Z_2 further broadcasts this traffic to
17 B. Z_1 and Z_2 include WPAN functionality and may act as data relay agents only, and thus
18 may not process data. Piconet 1 and piconet 2 include respective PNC₁ and PNC₂ and
19 thus devices A and B only need PNC₁ and PNC₂, respectively, for allocation of time
20 slots, and the function of protection of content is performed by the security manager 24
21 of each device A, B.

22 [0034] In order to facilitate communication between devices A and B, in different
23 piconets 1 and 2, device A is associated with a router 28 which stores information related
24 to other devices in its piconet 1, and routing information having instructions on how to
25 route traffic from device A to other devices, such as device B. Correspondingly, device B
26 is also associated with a router 30 having similar functionalities. Thus, any device A or B
27 is associated with a router and these routers 28, 30 query each other periodically in order
28 to update router information, due to the dynamic nature of the ad hoc networks 18.

29 [0035] Referring to Figure 7 and Figure 8, in order to establish a secure
30 communication between device A and B, device A performs the steps of acquiring device
31 B's full static address or device ID and a public key or symmetric key in order to perform

1 key agreement, in step 110. In the next step 112, the key agreement yields an
2 authentication key for subsequent communication. Once device A receives a response, in
3 predetermined time, that proves possession of the group public key, in step 114, then
4 device A generates a new set of group keys and transports these keys to device B, in step
5 116. Device B can then acknowledge receipt of group keys in step 118. Thus, devices A
6 and B require each other's authentic public key and each other's full device ID for
7 authentication and establishment of a secure channel 26, as different piconets may use
8 different short hand address addresses for each device A or B. Therefore, device A and
9 device B form a trusted group and a secure channel is set up if device B trusts any of the
10 intermediate routers, otherwise device B creates its own keys in order to set up a secure
11 channel 26

12 **[0036]** Although the invention has been described with reference to certain specific
13 embodiments, various modifications thereof will be apparent to those skilled in the art
14 without departing from the spirit and scope of the invention as outlined in the claims
15 appended hereto.

**THE EMBODIMENTS OF THE INVENTION IN WHICH AN EXCLUSIVE
PROPERTY OR PRIVILEGE IS CLAIMED ARE DEFINED AS FOLLOWS:**

1. A method of establishing and maintaining distributed security between a plurality of devices in an ad hoc network, the method having the steps of;

associating each device with a unique device address;

assigning to one of said devices a control function to control access by other devices to said network;

each of said devices generating a public key for distribution to other devices;

each of said devices authenticating itself periodically with said other devices in order to determine status of said other devices;

arranging said devices into a plurality of trust groups, each group having a group key for distribution within said trust group;

associating a trust level to each of said devices;

each of said devices using said public key and said group key to perform key agreement in order to establish a secure communication channel with said other devices in said group;

whereby each of said devices is responsible for its own security by generating, distributing its own keys to said other devices.

2. The method of claim 1 wherein said device determines a source of said group key.

3. The method of claim 2 wherein when said source is a device in said trust group then said group key is used for encryption and decryption of data transmitted between said devices.

4. The method of claim 2 wherein when said source is a device excluded trust group then said group key is used decryption of data transmitted to said device.

5. The method of claim 1 wherein step of determining status of said other devices includes a further step of determining which of said devices are active and capable of participating in said network.

6. The method of claim 1 wherein step of determining status of said other devices includes a further step of using a challenge response protocol using said group key to establish whether said other devices are allowed access to said network in accordance with said control function.

7. The method of claim 1 wherein said unique device address includes a device ID or a local ID.
8. The method of claim 7 wherein said device ID is an IEEE MAC address and said local ID is an n-bit address unique to said group.
9. A method of establishing and maintaining distributed security between one correspondent and another correspondent, said correspondents being members of different ad hoc networks and forming a group of communicating correspondents, the method having the steps of;
- associating said one correspondent and said other correspondent with a unique device address;
 - controlling access to said different ad hoc networks;
 - each ad hoc network having a gateway and transferring traffic between said correspondents via said gateways;
 - said one correspondent generating a public key for distribution to said other correspondent;
 - said one correspondent authenticating itself periodically with said other correspondent in order to determine status of said other correspondent;
 - determining a group key for distribution to said correspondents in accordance to said step of controlling access;
 - associating a trust level to each of said correspondents;
 - each of said correspondents using said public key and said group key for performing key agreement in order to establish secure communication within said group;
 - whereby each of said correspondents is responsible for its own security by generating, distributing its own keys to said other devices.
10. The method of claim 9 wherein said step of transferring traffic includes a further step of associating each of said correspondents with a router for storing routing information having instructions for routing traffic from said one correspondent to said other correspondent.
11. The method of claim 10 wherein said routers query each other periodically in order to update and maintain said routing information.

12. The method of claim 11 wherein said step of determining said status of said other correspondent includes a further step of using a challenge response protocol to establish whether said other correspondent is allowed access to said different ad hoc network having said one correspondent, in accordance with said control function.

13. A distributed security system for a plurality of devices in a communication network, each of said devices being responsible for generating, distributing and controlling its own keys for access to said communication network and using said keys to establish a trusted network, each device's membership to said communication network being checked periodically by other devices by using a challenge response protocol to establish which devices are allowed access to said communication network and said trusted network.

14. The system of claim 13 wherein each device includes a security manager having the functions of generating said keys and distributing said keys to selected devices in said trusted network.

15. The system of claim 14 wherein said trusted network is associated with a level of trust.

16. The system of claim 14 wherein said security manager determines a source of said keys such that said keys from a device within said trusted network may be used for encryption and decryption of data, and said keys from a device excluded from said trusted network may be used decryption of said data.

17. The system of claim 16 wherein said security manager foregoes decrypting said data when said keys are from a device excluded from said trusted network.

18. The system of claim 15 wherein an outcome of said periodic checking is recorded by said security manager in order to maintain and update a membership list, and adjust said level of trust accordingly.

19. The system of claim 17 wherein different trusted networks may be established within said network based on differing levels of trust.

20. The system of claim 13 wherein said communication network includes a plurality of ad hoc networks and said distributed security system is established between devices in different ad hoc networks.

21. The system of claim 19 wherein each ad hoc network includes a controller to controlling access to each of said ad hoc networks, each ad hoc network having a gateway for transferring traffic therebetween, and device having a router for storing routing information having instructions for routing traffic from said one device to another device via said gateways and other routers.

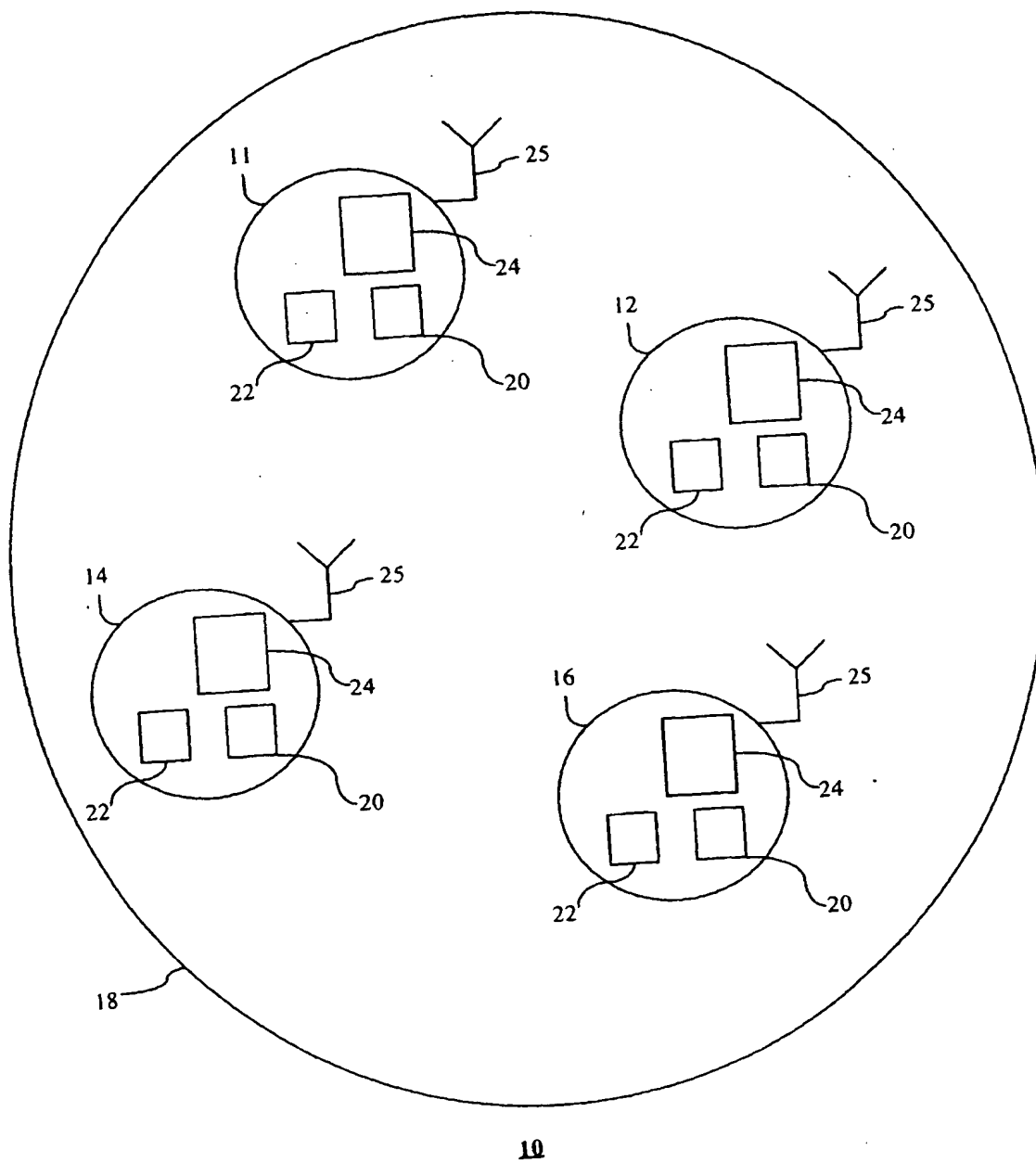


Figure 1

	A B C D E F G H		
Group 1'	x x x	Key source: C	encryption/decryption permitted
Group2'	x x	Key Source: D	decryption only

Figure 2

	A B C D E F G H		
Group 1'	x x x	Key source: C	encryption/decryption
Group2'	x x	Key Source: D	decryption
Group3'	x x	Key source: A	encryption/decryption

Figure 3

	A B C D E F G H		A	D
Group 1'	x x x	Key source: C	encryption/decryption	
Group2'	x x x	Key Source: G	encryption/decryption	decryption
Group3'	x x	Key Source: E		decryption

Figure 4

	A B C D E F G H		A	D
Group 1	x x x	Key source: C	encryption/decryption	
Group2	x x \$ x	Key Source: G	encryption/decryption	decryption
Group3'	x x	Key Source: E		decryption

\$: hidden node ('fly on the wall')

Figure 5

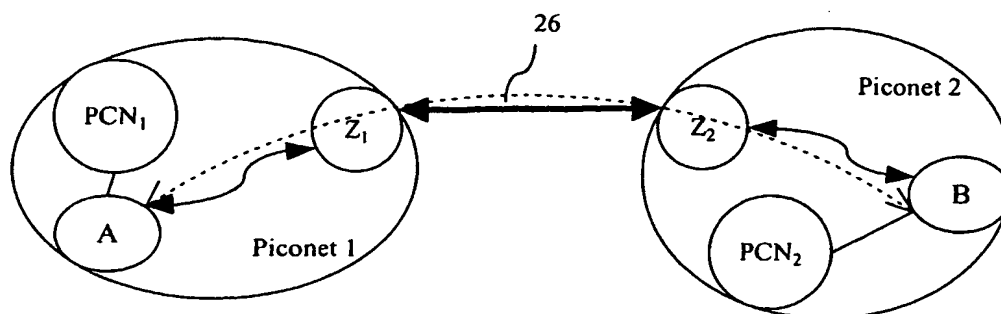


Figure 6

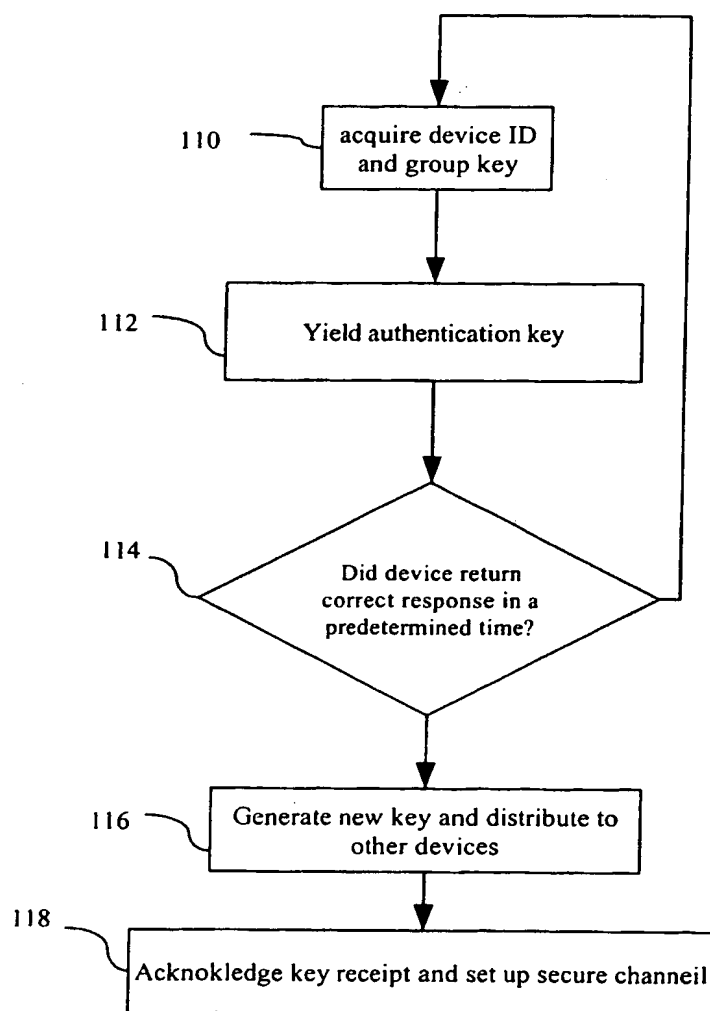


Figure 7

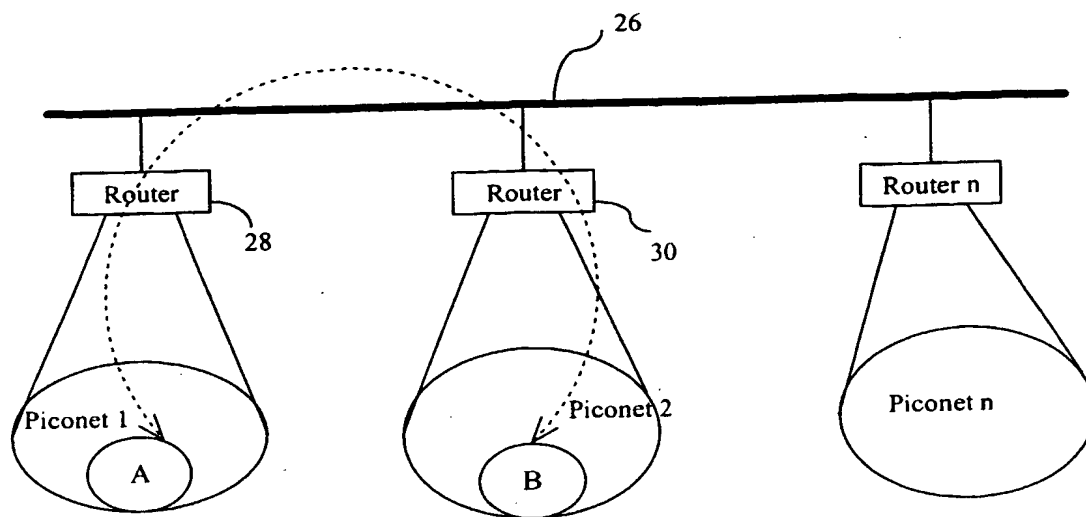


Figure 8

INTERNATIONAL SEARCH REPORT

International Application No

PCT/CA 03/00315

A. CLASSIFICATION OF SUBJECT MATTER

IPC 7 H04L29/06

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 7 H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

EPO-Internal

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X A	<p>VENKATRAMAN L ET AL: "A novel authentication scheme for ad hoc networks" DEPARTMENT OF ELECTRICAL AND COMPUTER ENGINEERING AND COMPUTER SCIENCE, vol. 3, 23 September 2000 (2000-09-23), pages 1268-1273, XP010532729</p> <p>page 1269, right-hand column, line 28 -page 1271, left-hand column, line 13 page 1272, left-hand column, line 15 -page 1273, left-hand column, line 11</p> <p style="text-align: center;">--- -/--</p>	<p>1-3, 5-7, 9-15, 20, 21</p> <p>4, 8, 16-19</p>

☒ Further documents are listed in the continuation of box C.☒ Patent family members are listed in annex.

* Special categories of cited documents:

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S document member of the same patent family

Date of the actual completion of the international search

9 July 2003

Date of mailing of the international search report

16/07/2003

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Lázaro, M.L.

International Application No
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PCT/CA 03/00315

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	JACOBS S ET AL: "MANET Authentication Architecture" INTERNET DRAFTS, March 1999 (1999-03), XP002206450 Retrieved from the Internet: <URL:http://www.watersprings.org/pub/id/draft-jacobs-imep-auth-arch-00.txt> 'retrieved on 2002-07-17!	13,14, 20,21
A	page 2, line 19-32 page 6, line 5-28 page 7, line 13 -page 8, line 15 page 13, line 19 -page 14, line 8 page 15, line 1-25 --- EP 1 102 430 A (ERICSSON TELEFON AB L M) 23 May 2001 (2001-05-23)	1-12, 15-19
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International Application No

PCT/CA 03/00315

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